

## Expressing languages with logic

The set of words of the form  $a^n b^n$  can be defined with:

$$\exists x, (\forall y, y \leq x \Rightarrow a(y)) \wedge (\forall y, y > x \Rightarrow b(y)) \wedge (\exists z, \text{End}(z) \wedge \text{Half}(x, z))$$

The expressivity of a logic can depends on:

Possible quantification allowed:

- The full first-order logic:  $\text{FO}$ ,
- Restricting to two variables:  $\text{FO}_2$ ,
- Adding modular quantifiers:  $\text{FO} + \text{MOD}$ ,
- Bounded alternation of quantifiers:  $\Sigma_k$ ,
- Many more...

Possible numerical predicates:

- The order:  $<$ ,
- The regular predicates:  $+1$ ,
- The modular predicates:  $\text{MOD}$ ,
- Any arbitrary predicates:  $\text{ARB}$ ,
- Many more...

$\left. \begin{array}{l} \text{The order: } <, \\ \text{The regular predicates: } +1, \\ \text{The modular predicates: } \text{MOD}, \\ \text{Any arbitrary predicates: } \text{ARB}, \end{array} \right\} \text{The regular predicates: } \text{REG}.$

## Straubing's central conjecture

Let  $\mathcal{L}$  be a logic,

$$\mathcal{L}[\text{ARB}] \cap \text{Reg} = \mathcal{L}[\text{REG}] ?$$

The regular languages.

For which logics  $\mathcal{L}$  does the regular languages in  $\mathcal{L}[\text{ARB}]$  are exactly the languages of  $\mathcal{L}[\text{REG}]$ ?

True

- $\Sigma_1$ ,
- $\text{FO}$ ,
- $\text{FO} + \text{MOD}(p^k)$ , for a fixed prime  $p$ .

False

- (the exotic)  $\text{FO} + S_5$ .

Open

- $\text{FO} + \text{MOD}$ ,
- $\text{FO}_2$ ,
- $\Sigma_k$ .

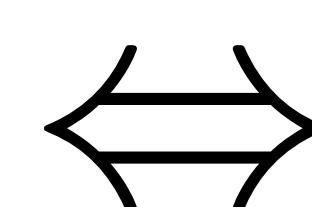
## The logic $\Sigma_2[\text{ARB}]$

$$\exists x_1, \dots, x_k \quad \forall y_1, \dots, y_k \quad \varphi(x_1, \dots, y_k)$$

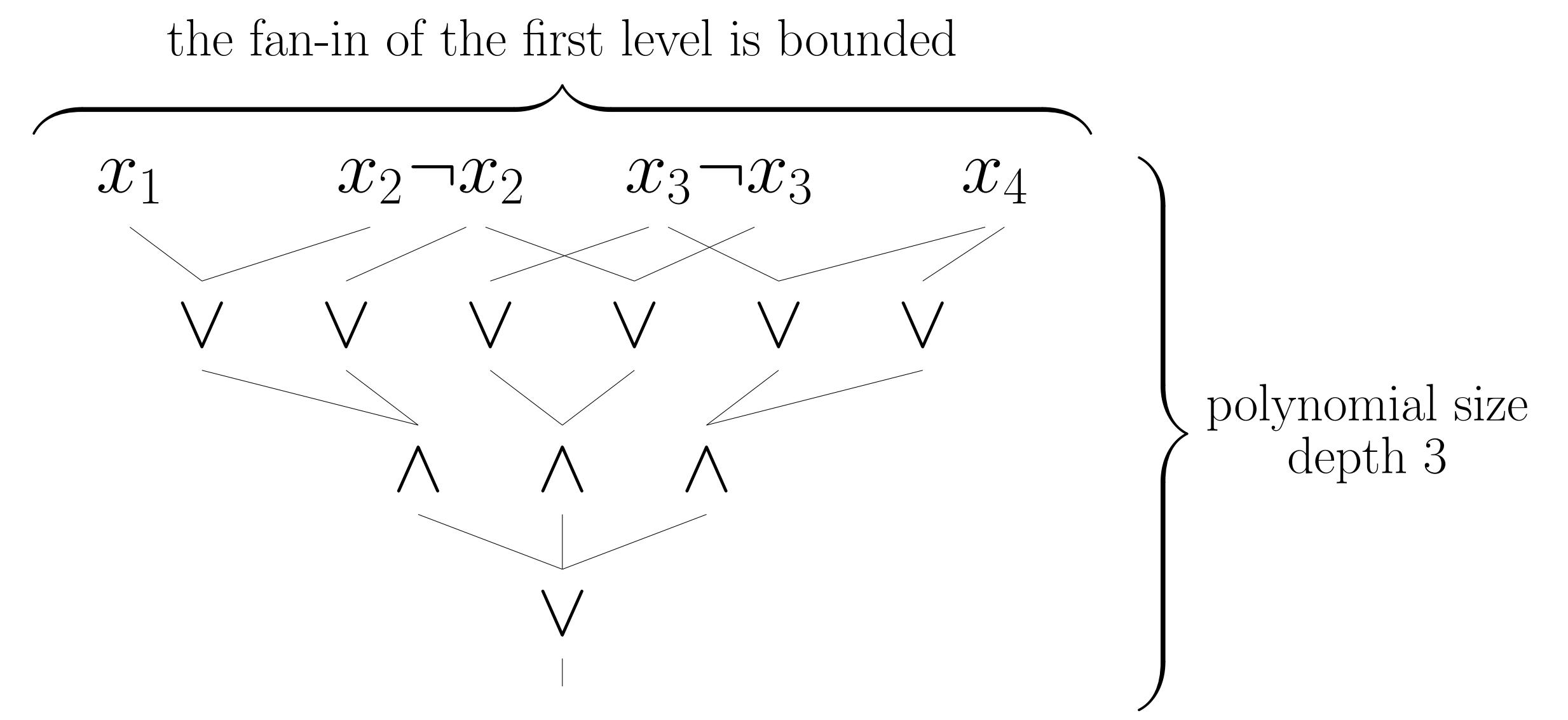
A block of existential quantifiers

A block of universal quantifiers

A formula that can use arbitrary predicates



## The circuit class $\exists \forall \exists$



## The main result (with Cadilhac, Paperman and Zeume)

$$\Sigma_2[\text{ARB}] \cap \text{Reg} \cap \text{Neut} = \Sigma_2[\text{REG}] \cap \text{Neut} .$$

The two parts of the proof:

- An algebraic characterisation of  $\Sigma_2[<]$  (Pin and Weil).
- Lower bounds against  $\exists \forall \exists$ .

The class of languages with a neutral letter: a mild technical assumption.

## A corollary of the proof

The logic  $\Pi_2$  is defined as  $\Sigma_2$  but with an initial block of universal quantifiers.

The logic  $\Delta_2$  is defined as the class of formulas that can be written both as a  $\Sigma_2$  formula and as a  $\Pi_2$  formula.

$$\Delta_2[\text{ARB}] \cap \text{Reg} = \Delta_2[\text{REG}] .$$