

# THE REGULAR LANGUAGES OF FIRST-ORDER LOGIC WITH ONE ALTERNATION



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## Expressing languages with logic

The set of words of the form  $a^n b^n$  can be defined with:

$$\exists x, (\forall y, y \leq x \Rightarrow a(y)) \wedge (\forall y, y > x \Rightarrow b(y)) \wedge (\exists z, \text{End}(z) \wedge \text{Half}(x, z))$$

The expressivity of a logic can depends on:

Possible quantification allowed:

- The full first-order logic: **FO**,
- Restricting to two variables: **FO<sub>2</sub>**,
- Adding modular quantifiers: **FO + MOD**,
- Bounded alternation of quantifiers: **Σ<sub>k</sub>**,
- Many more...

Possible numerical predicates:

- The order: **<**,
  - The regular predicates: **+1**,
  - The modular predicates: **MOD**,
  - Any arbitrary predicates: **ARB**,
  - Many more...
- } The regular predicates: **REG**.

## Straubing's central conjecture

Let **ℒ** be a logic,

$$\mathcal{L}[\text{ARB}] \cap \text{Reg} = \mathcal{L}[\text{REG}] \text{ ?}$$

The regular languages.

For which logics **ℒ** does the regular languages in **ℒ[ARB]** are exactly the languages of **ℒ[REG]**?

True

- **Σ<sub>1</sub>**,
- **FO**,
- **FO + MOD(*p<sup>k</sup>*)**, for a fixed prime *p*.

False

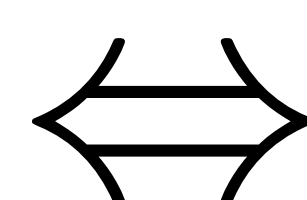
- (the exotic) **FO + S<sub>5</sub>**.

Open

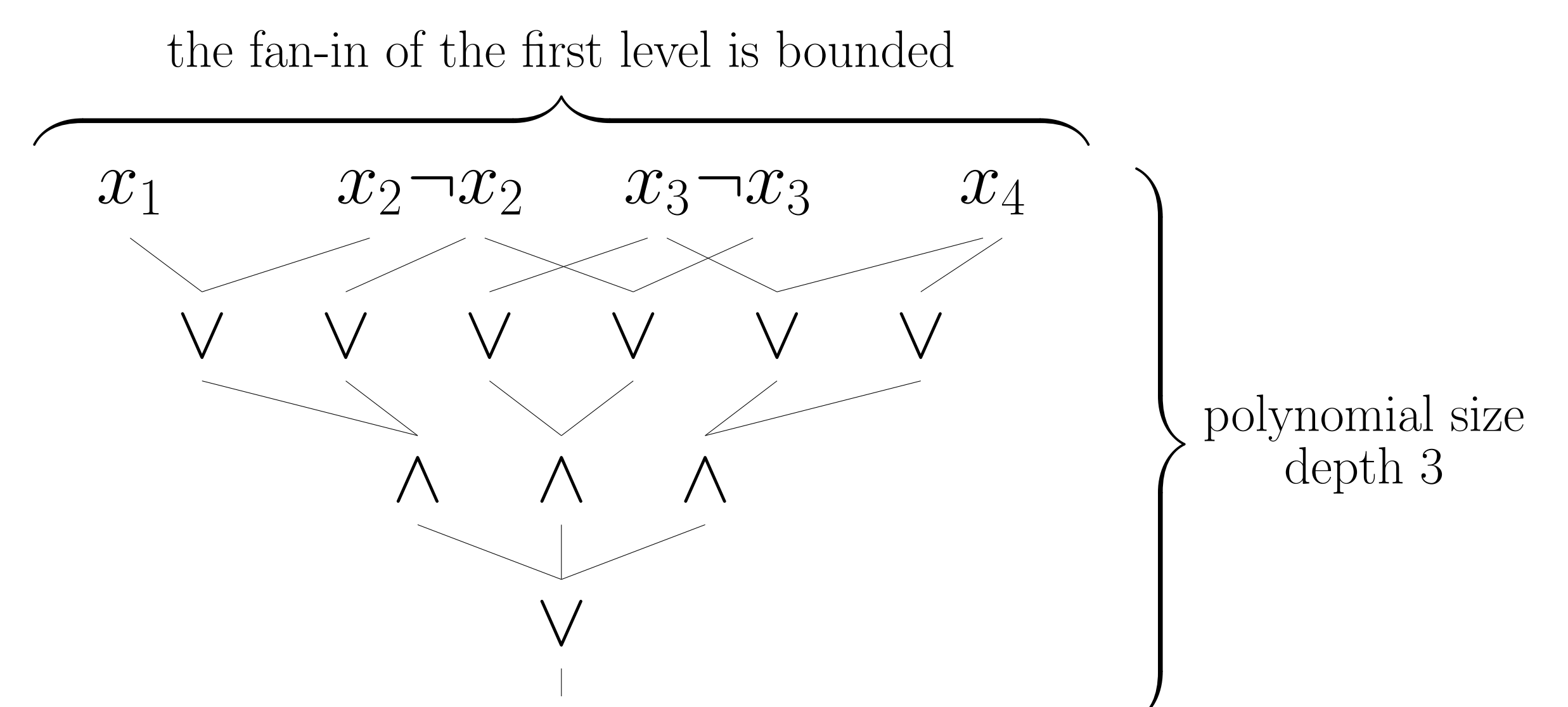
- **FO + MOD**,
- **FO<sub>2</sub>**,
- **Σ<sub>k</sub>**.

## The logic Σ<sub>2</sub>[ARB]

$$\underbrace{\exists x_1, \dots, x_k}_{\text{A block of existential quantifiers}} \underbrace{\forall y_1, \dots, y_k}_{\text{A block of universal quantifiers}} \underbrace{\varphi(x_1, \dots, y_k)}_{\text{A formula that can use arbitrary predicates}}$$



## The circuit class ∃∀∃



## The main result (with Cadilhac, Paperman and Zeume)

$$\Sigma_2[\text{ARB}] \cap \text{Reg} \cap \text{Neut} = \Sigma_2[\text{REG}] \cap \text{Neut} .$$

The two parts of the proof:

- An algebraic characterisation of **Σ<sub>2</sub>[<]** (Pin and Weil).
- Lower bounds against ∃∀∃.

The class of languages with a neutral letter: a mild technical assumption.

## A corollary of the proof

The logic **Π<sub>2</sub>** is defined as **Σ<sub>2</sub>** but with an initial block of universal quantifiers.

The logic **Δ<sub>2</sub>** is defined as the class of formulas that can be written both as a **Σ<sub>2</sub>** formula and as a **Π<sub>2</sub>** formula.

$$\Delta_2[\text{ARB}] \cap \text{Reg} = \Delta_2[\text{REG}] .$$